## THEORY OF COMPUTATION (CSEN 5234)

Time Allotted : 3 hrs

1.

## Figures out of the right margin indicate full marks.

## Candidates are required to answer Group A and any 5 (five) from Group B to E, taking at least one from each group.

Candidates are required to give answer in their own words as far as practicable.

## Group - A (Multiple Choice Type Questions)

- Choose the correct alternative for the following: (i) A minimum state DFA accepting the language L = {w |  $w \in (0,1)^*$ , where, the number of 0's and 1's of w are divisible by 5 and 4, respectively} has
  - (b) 20 states (c) 15 states (a) 10 states (d) 16 states
  - The Pumping Lemma for Regular Languages can be used to prove that a given (ii) language (a) is Type 2 (b) is Type 3 (c) is not Type 2 (d) is not Type 3.
- Let S & T be language over  $\Sigma$ : {a, b} represented by the following regular expressions (iii)  $(a + b^*)^*$  and  $(a + b)^*$  respectively, then which of the following is TRUE? (a)  $S \subset T$ (b)  $T \subset S$ (c) S = T(d)  $S \cap T = \emptyset$ .
- (iv) Consider the following three languages:  $L_1 = \{ww^R \mid w \in (0,1)^*\}$  $L_2 = \{w \# w^R \mid w \in (0,1)^*\}, where \# is a special symbol$  $L_3 = \{ww \mid w \in (0,1)^*\}$ Which one of the following is false? (a) L1 is a deterministic CFL (b) L2 is a deterministic CFL (c) L3 is a CFL (d) both (a) and (c).
- (v) Suppose that L<sub>4</sub> and L<sub>5</sub> are two languages (over the same alphabet) given to you such that both  $L_4$  and  $L_4 \cap L_5$  are CFL. Then which of the following is incorrect? (a)  $L_5$  must be CFL too (b)  $L_5$  can never be regular (c)  $L_5$  can be regular (d) Both (a) & (b).
- (vi) If the original nondeterministic finite state automata (ndfsa) has n states, the corresponding deterministic finite state automata (dfsa) cannot have (a) more than n states (b) more than 2n states
  - (c) more than  $2^n$  states

(d) more than (n-1) states.

 $10 \times 1 = 10$ 

Full Marks: 70

- Which one of the following languages cannot be accepted by a non-deterministic (vii) pushdown acceptor (ndpda)?
  - (a) {  $0^m 1^n | 0 < m, 0 < n$  } (c) {  $0^{m}1^{m}0^{m} | 0 < m$  }
- (b) {  $0^m 1^n 0^m | 0 < m, 0 < n$  }

(d) {  $0^m 1^n 0^r | 0 < m, 0 < n, 0 < r$  }.

### If a Context Free Grammar is in Chomsky Normal Form then a production (viii)

- (a) may contain a single non-terminal in the RHS
- (b) may contain both terminal and non-terminal symbols in the RHS
- (c) may be of the form  $A \rightarrow \varepsilon$  where A is not the start symbol
- (d) none of these.
- If a Context Free Language L<sub>3</sub> has a Type 2 grammar G which is ambiguous then (ix) there must be a terminal string  $\alpha$  such that
  - (a)  $\alpha$  has two different derivation trees
  - (b)  $\alpha$  has two different rightmost derivations
  - (c)  $\alpha$  has two different leftmost derivations
  - (d) all of the above.
- Let L<sub>6</sub> be a recursive language, and L<sub>7</sub> be a recursively enumerable language but (x) not recursive language. Which of the following is true?
  - (a)  $L_6'$  is recursive  $L_7'$  is recursively enumerable
  - (b)  $L_6'$  is recursive  $L_7'$  is not recursively enumerable
  - (c)  $L_6'$  and  $L_7'$  are recursively enumerable
  - (d)  $L_6'$  is recursively enumerable  $L_7'$  is recursive.

# Group – B

A non-deterministic finite state acceptor (ndfsa) M<sub>2</sub> has the state table shown 3. (a) below. S is the start state, and C is the only goal state. Convert  $M_2$  to an equivalent deterministic finite state acceptor (dfsa) M<sub>3</sub> and clearly indicating the start and goal states.

	0	1	
→S	Р	Q	
Р	Р, А	B,C	
Q	Q, A	B,C	
Α	A,C		
В		С	
*C	С	С	
DFA	over alph	abet {0, 1}	that accepts thos

[(CO1)(Apply/IOCQ)]

5 + 5 + 2 = 12

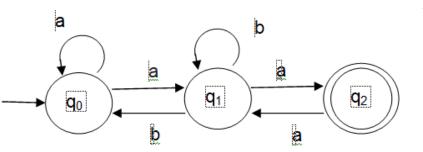
- (b) Construct an DFA over alphabet {0, 1} that accepts those strings that contain the [(CO2)(Create/HOCQ)] pattern 001 somewhere. [(CO1)(Remember/LOCQ)]
- What do you mean by 2- equivalent states? (c)
- 2. Design a DFA that accepts every string over  $\Sigma = \{a, b\}$  which starts and ends in (a) same symbol. [(CO1)(Create/IOCQ)]
  - Consider the DFA given below (q1 is the starting state and q3, q5 are two final (b) states)

δ	0	1
$\mathbf{q}_1$	<b>q</b> <sub>2</sub>	<b>q</b> <sub>3</sub>
<b>q</b> <sub>2</sub>	<b>q</b> 3	<b>q</b> 5
<b>q</b> 3	<b>q</b> 4	<b>q</b> 3
<b>q</b> 4	<b>q</b> 3	<b>q</b> 5
<b>q</b> 5	<b>q</b> 2	<b>q</b> 5
1	1 1 1	

Find out the indistinguishable and distinguishable states for the automata.Construct minimum state equivalent of automata.[(C01)(Apply/L0CQ)]6 + 6 = 12

## Group – C

- 4. (a) Consider the language  $L_8 = \{x \in \{0, 1\}^* \mid x \text{ does not start with } 010\}$ . Show that the language is regular. [(CO2)(Apply/IOCQ)]
  - (b) Let L<sub>10</sub> be the set of all palindromes over {a, b}. Construct a grammar G generating L<sub>10</sub>. According to Chomsky classification, which type of grammar is this? Can we construct a nondeterministic finite automaton for this language? Justify your answer.
    - 5 + (3 + 2 + 2) = 12
- 5. (a) Prove that  $(1 + 00^*1) + (1 + 00^*1)(0 + 10^*1)^* (0 + 10^*1) = 0^*1(0 + 10^*1)^*$ [(CO2)(Apply/IOCQ)]
  - (b) Consider the transition system given in the figure, below.
     Prove that the strings recognized are given by (a + a(b + aa)\*b)\* a(b + aa)\* a.
     [(CO2)(Apply/IOCQ)]



(c) What do you mean by right linear grammar? Give an example. (CO6)(Under CO6)(Under CO6)(Under

[(CO6)(Understand/LOCQ)]3 + 6 + 3 = 12

## Group - D

6. (a) Simplify the given CFG.  $S \rightarrow ACD$   $A \rightarrow a \mid F$   $B \rightarrow \varepsilon$   $C \rightarrow ED \mid \varepsilon$   $D \rightarrow BC \mid b$  $E \rightarrow b$ 

$$F \rightarrow aF$$

(b) Consider the CFG given below:  $S \rightarrow a \mid Sa \mid bSS \mid SSb \mid SbS$ Is the CFG ambiguous? Justify your answer.

[(CO3)(Apply/IOCQ)] 6 + 6 = 12

- 7. (a) Design a pushdown automata accepting  $L= \{w \in \{a,b\}^* : number of a's in w is exactly double of the number of b's in w\}.$  [(CO4)(Create/HOCQ)]
  - (b) Using the Pumping Lemma for Context-Free Languages, show that the language  $L = \{0^a 1^b 0^a 1^b | a > 0, b > 0\}$  is not a Type 2 (context-free) language. [(CO3)(Evaluate/HOCQ)] 6 + 6 = 12

## **Group – E**

8. (a) Nobody knows yet if P = NP. Consider the language L<sub>11</sub> defined as follows:  $(0 + 1)^*$  if P = NP.

$$L_{11} = \begin{cases} (0 + 1)^{-1} & \text{if } \\ 0 & \text{otherwise} \end{cases}$$

Is L<sub>11</sub> recursive or recursively enumerable? Justify your answer. [(CO7)(Analyse/IOCQ)]

(b) A single tape Turing machine (M) has three states q0, q1, and q2, where, q0 is the starting state. The tape symbols of M are {a, b, B} and I/P alphabets are {a, b}. The symbol B is the blank symbol. The transition function of the M is given below:

	а	b	В
q0	q1, a, R	-	q0, B, R
q1	-	q2, b, R	-
q2	q2, a, R	q2, b, R	Halt

What is the language accepted by the Turing machine M?[(C05)(Apply/LOCQ)]Write a short note on multi-tape Turing machine.[(C05)(Remember/LOCQ)]

4 + 4 + 4 = 12

9. (a) Design a Turing machine that can compute the sum of two numbers. [(CO5)(Apply/IOCQ)]
(b) Design a Turing machine M<sub>11</sub> that find the 1's complement of a binary number. e.g. if the input number is 00110 then the output will be 11001. [(CO5)(Create/HOCQ)]

6 + 6 = 12

Cognition Level	LOCQ	IOCQ	HOCQ
Percentage distribution	26.04	36.46	37.50

### Course Outcome (CO):

(c)

After the completion of the course students will be able to

- 1. Design and analyze Deterministic and non-deterministic finite state automata.
- 2. Understand the correspondence between finite state automata and regular languages.
- 3. Design context free grammars to generate strings from a context free language and convert them into Chomsky normal forms.
- 4. Design deterministic and non-deterministic push down automata to recognize context free languages.
- 5. Construct Touring machines for computable functions.
- 6. Understand the hierarchy of formal languages, grammars and machines.
- 7. Distinguish between computability and non-computability and Decidability and undecidability.

\*LOCQ: Lower Order Cognitive Question; IOCQ: Intermediate Order Cognitive Question; HOCQ: Higher Order Cognitive Question