



- (vii) A deterministic pushdown acceptor (dpda) can be constructed to recognize any given language L of Type n, where n equals  
 (a) 0                      (b) 1                      (c) 2                      (d) 3.
- (viii) The Pumping Lemma for Regular Languages can be used to prove that a given language L  
 (a) is not Type 0    (b) is not Type 1    (c) is not Type 2    (d) is not Type 3.
- (ix) A machine M has been designed such that it accepts a given positive integer n as input, if and only if n is a prime number. Then which one of the following alternatives is false?  
 (a) M cannot be a deterministic finite state acceptor (dfsfa)  
 (b) M cannot be a non-deterministic finite state acceptor (ndfsa)  
 (c) M cannot be a non-deterministic pushdown acceptor (ndpda)  
 (d) M cannot be a Turing machine.
- (x) We are supplied two deterministic finite state acceptors (dfsas)  $M_1$  and  $M_2$  on the input alphabet  $\{0,1\}$ . We want to determine whether the languages  $L(M_1)$  and  $L(M_2)$  are identical sets, i.e., whether they contain exactly the same strings. Which of the following alternatives is true?  
 (a) There is a dfsa that will solve the problem  
 (b) There is a pda that will solve the problem but not a dfsa  
 (c) There is a Turing machine that will solve the problem but not a pda  
 (d) There is no effective procedure for solving the problem.

**Group - B**

- 2.(a) Design a deterministic finite state acceptor (dfsfa)  $M_1$  that will accept only those strings on the alphabet  $\{0,1\}$  which do *not* begin with the substring '010'. Show both the state table and the state transition diagram of  $M_1$  and briefly explain how  $M_1$  works.
  - (b) Design a deterministic finite state acceptor (dfsfa)  $M_2$  that will accept only those strings on the alphabet  $\{0,1,2\}$  in which the symbol 2 never follows the symbol 0. For example,  $M_2$  will accept the strings 101200 and 10101 but not the string 110020. Show both the state table and the state transition diagram of  $M_2$ .
- 6+6=12**
- 3.(a) A non-deterministic finite state acceptor (ndfsa)  $M_3$  has the state table shown below. The start state is S and the only final state is C. Convert  $M_3$  to an equivalent deterministic finite state acceptor (dfsfa)  $M_4$ , clearly indicating the start and final states. Briefly explain your method of conversion.

	0	1
→ S	S, A	S
A	--	B
B	--	C
*C	C	C

- (b) Minimize the number of states in the machine  $M_4$  designed above to get a dfsa  $M_5$ . Briefly explain the method of minimization.

**6+6=12**

**Group - C**

- 4.(a) Construct a deterministic finite state acceptor (dfsfa)  $M_6$  on the input alphabet  $\{0,1\}$  that accepts a string  $\alpha$  if and only if  $\alpha$  is contained in the regular expression  $0^*(0+1)(0+1)^*$ .

- (b) Construct a deterministic finite state acceptor (dfsfa)  $M_7$  on the input alphabet  $\{0,1\}$  that accepts a string  $\alpha$  if and only if  $\alpha$  is *not* contained in the regular expression  $0^*1^*01$ .

**6+6=12**

- 5.(a) State and explain the Pumping Lemma for Regular Languages. Is it possible to use this lemma to show that a given language  $L$  is *not* context-free?

- (b) Use the Pumping Lemma for Regular Languages to show that the language  $L=\{0^m1^n0^m1^n \mid m > 5, n > 10\}$  is not regular.

**6+6=12**

**Group - D**

- 6.(a) When is a Type 2 (context-free) grammar said to be ambiguous? Consider the grammar  $G = (V, T, S, P)$  where  $V = \{S, A, B, 0, 1, 2, 3, 4\}$ ,  $T = \{0, 1, 2, 3, 4\}$ ,  $S$  is the start symbol, and the productions are as follows:  $S \rightarrow A$ ,  $S \rightarrow B$ ,  $S \rightarrow 4$ ,  $A \rightarrow 012S$ ,  $B \rightarrow 012S3S$ . Show that the grammar  $G$  is ambiguous.

- (b) Provide a Type 2 (context-free) grammar for the following language  $L$  over the input alphabet  $\{0, 1\}$ :  $L = \{\alpha \mid \text{the string } \alpha \text{ has more 0's than 1's}\}$

**6+6=12**

- 7.(a) Let  $G$  be a Type 2 (context-free) grammar in Chomsky Normal Form. Let  $\alpha$  be a string in  $L(G)$  of length  $n$ , i.e.,  $\alpha$  has  $n$  terminal symbols. Explain why the parse tree of  $\alpha$  must have exactly  $(2n - 1)$  internal nodes (i.e., nodes with children).

- (b) Using the Pumping Lemma for Context-Free Languages, show that the language  $L=\{0^a1^b0^a1^b \mid a > 0, b > 0\}$  is not a Type 2 (context-free) language.

**6+6=12**

**Group - E**

8.(a) Two positive integers  $m$  and  $n$  are written on a Turing machine tape in unary notation. The integer  $m$  is to the left of the integer  $n$  on the tape, and a single blank cell separates the two integers. At start the read/write head is positioned on the leftmost 1 of  $m$ . Give the state transition diagram of a Turing machine that will halt on the leftmost 1 of  $m$  if  $m \geq n$ , and will halt on the leftmost 1 of  $n$  if  $m < n$ . Clearly state any assumptions made.

(b) When is a set of integers said to be a recursive set? What are the major characteristics of a recursive set? Explain whether the set of all positive integers that are not prime numbers is a recursive set.

**6+6=12**

9.(a) Briefly clarify what is meant by a 'Universal Turing Machine'. Can every effective procedure be implemented by a Universal Turing Machine assuming the input is appropriately supplied?

(b) What is the 'Halting Problem' for Turing Machines? "Halting Problem for Turing machines is unsolvable" - comment.

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